

## Worksheet 9. Roots of unity and polynomial multiplication

**Euler's Formula.**  $e^{i\theta} = \cos\theta + i\sin\theta$  (which remember represents the point  $(\cos\theta, \sin\theta)$  on the unit circle in  $\mathbb{C}$ .)

**Definition.** An  $n$ th root of unity is a solution (in  $\mathbb{C}$ ) to  $z^n = 1$ .

**Problem 1.**

- (a) Prove that for any integer  $k$ , the number  $e^{2\pi ik/n}$  is a complex  $n^{\text{th}}$  root of unity. Where does it appear on the unit circle?  
 (b) Find all solutions  $\theta \in \mathbb{R}$  to  $e^{i\theta} = 1$ .  
 (c) Prove that, for any  $n \in \mathbb{N}$ , the numbers

$$e^{2\pi ik/n}, \quad k = 0, 1, \dots, n-1,$$

are the complex  $n^{\text{th}}$  roots of unity. (In particular, you must show that this is a list of  $n$  distinct numbers!) Draw a picture and indicate where these  $n$  points appear in the plane.

- (d) Write  $\zeta = e^{2\pi i/n}$ . Prove that

$$1, \zeta, \zeta^2, \dots, \zeta^{n-1}$$

is also a complete list of the  $n^{\text{th}}$  roots of unity.

- (e) Prove that if  $n$  is even, then squaring the  $n^{\text{th}}$  roots of unity gives a list (with repetitions) of the  $(n/2)^{\text{th}}$  roots of unity.  
 (f) Prove that if  $n$  is even, then the  $n^{\text{th}}$  roots of unity come in  $\pm$  pairs:  $\xi$  is an  $n^{\text{th}}$  root of unity iff  $-\xi$  is. What about when  $n$  is odd?

**Polynomial multiplication** Given two polynomials  $A(x) = a_0 + a_1x + a_2x^2 + \dots + a_nx^n$  and  $B(x) = b_0 + b_1x + b_2x^2 + \dots + b_nx^n$ , we would like to compute the coefficients of the product

$$\begin{aligned} A(x)B(x) &= a_0b_0 + (a_0b_1 + a_1b_0)x + \dots + a_nb_nx^{2n} \\ &= c_0 + c_1x + \dots + c_{2n}x^{2n}. \end{aligned}$$

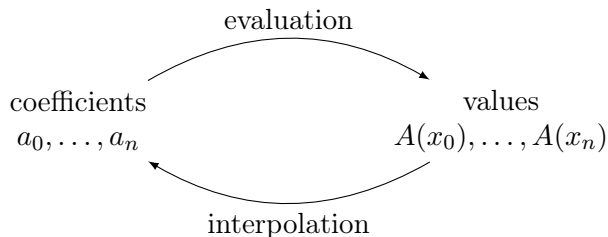
**Problem 2.** Find an explicit formula for the coefficient  $c_k$  of  $x^k$  in  $A(x)B(x)$ , for  $k = 0, 1, \dots, 2n$ .

**Problem 3.** Briefly discuss with your groupmates a naïve algorithm to multiply two degree  $n$  polynomials in  $O(n^2)$  time.

Our goal is to find a D&C solution that runs in  $O(n \log n)$  time. The main idea is to convert the polynomial to **point-value form**.

**Problem 4.** Discuss with your groupmates the assertion, *a polynomial of degree  $n$  is determined by  $n + 1$  of its values*. Can you interpret this in terms of linear algebra?

So we need to translate between coefficient form and point-value form efficiently:



**Problem 5.** Show how evaluation at a single value  $x$  can be performed in linear time using *Horner's Rule*:

$$A(x) = a_0 + x(a_1 + x(a_2 + \cdots + x(a_{n-2} + x(a_{n-1} + a_n x)) \cdots))$$

Do a small example, say a degree-3 polynomial.

So we need a way to interpolate quickly. The trick will be to choose the interpolation points  $x_k$  cleverly. But actually we won't worry much about interpolation yet; it will turn out by some magic that if we find a nice evaluation algorithm, then interpolation will fall right out of it.

**Problem 6.** Explain how, if we could both interpolate polynomials in  $O(n \log n)$  and evaluate at  $n$  points in  $O(n \log n)$  time, then we could multiply polynomials in  $O(n \log n)$  time. Draw a diagram.

**A preview:** Choose the  $n$  points for interpolation in  $\pm$  pairs, so that the even powers of  $\pm x_k$  are the same:

$$\pm x_0, \pm x_1, \dots, \pm x_{n/2-1}.$$

Then we can split  $A(x)$  up as a sum  $A(x) = A_E(x^2) + xA_O(x^2)$ , where  $A_E$  and  $A_O$  are each polynomials of degree  $\frac{n}{2} - 1$ . These lower-degree polynomials have to be evaluated at  $n/2$  points each:

$$(x_0)^2, (x_1)^2, \dots, (x_{n/2-1})^2.$$

But, (uh-oh!), these  $n/2$  points no longer come in  $\pm$  pairs! How do we continue the recursion?! **Answer:** By evaluating at the  $n^{\text{th}}$  roots of unity in  $\mathbb{C}$  (!), which we will explore on the next worksheet.

**In case you're fast, like last time:**

We want to interpolate! That is, we still want to be able to take  $n$  values of a polynomial  $A(x_0), A(x_1), \dots, A(x_{n-1})$  and return its coefficients  $a_0, a_1, \dots, a_{n-1}$ . This problem can be thought of in terms of matrices:

$$\begin{bmatrix} A(x_0) \\ A(x_1) \\ \vdots \\ A(x_{n-1}) \end{bmatrix} = \begin{bmatrix} 1 & x_0 & x_0^2 & \cdots & x_0^{n-1} \\ 1 & x_1 & x_1^2 & \cdots & x_1^{n-1} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & x_{n-1} & x_{n-1}^2 & \cdots & x_{n-1}^{n-1} \end{bmatrix} \begin{bmatrix} a_0 \\ a_1 \\ \vdots \\ a_{n-1} \end{bmatrix}.$$

**Problem 7** (Challenging!). The large  $n \times n$  matrix  $M$  is called a **Vandermonde matrix**. Prove that if the  $x_i$ s are distinct, then the Vandermonde matrix is invertible.